LECTURE 15 - Stochastic Bandits

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I - Introduction

We call stochastic (multi-armed) bandit problem the following sequential decision problem.

(In American slang, a "one-armed bandit" refers to a slot machine. A multi-armed bandit therefore refers to a collection of many slot machines)

First let N>2 be known to the player. Consider probability distributions $v_1,...,v_N$ over R, unknown to the player.

At each round $t \ge 1$, the player chooses an action $I_t \in \{1,...,N\}$ (in the bandit jargon, actions 1,...,N are also called arms an when the player chooses action I_t we say he "pulls" arm I_t).

At the same time, the environment

draws a random vector $l_t = (l_t(i), ..., l_t(N))$ where • $l_t(i) \sim v_i$

• Variables $\{l_t(i)\}_{i=1}^N$ are independent t• Variable $l_t(i)$ is independent from $\{I_s\}_{s=1}^s$

and $\{l_s(j)\}_{1\leqslant s\leqslant t-1,1\leqslant j\leqslant N}$ The environment only reveals l_{I_t} , t to the player. He incurs this loss and moves on to the next round.

General Goal:

Minimize the regret which is defined here as

$$R_T = \mathbb{E}\left[\sum_{t=1}^T l_t(I_t)\right] - \min_{1 \leq i \leq N} \mathbb{E}\left[\sum_{t=1}^T l_t(i)\right].$$

Alternative representation of the regret:

Notation)

Mi: mean value of probability measure si

 $m^* = min m_i$ and $i^* \in argmin m_i$ $1 \le i \le N$ (in particular $m^* = m_{i^*}$)

in parnum. $\Delta_{i} = m_{i} - m^{*}$ t $n_{i}(t) = \sum_{s=1}^{s} 11 \{T_{s} = i\} : \text{Number of times}$ the player took action i up to time t.

Lemma 1

$$R_{T} = \sum_{i:\Delta_{i}>0} \Delta_{i} \mathbb{E} \left[n_{i}(T) \right]$$

Proof. Since
$$l_{t}(I_{t}) = \sum_{i=1}^{N} l_{t}(i) \, 1 \{I_{t} = i\}$$

we obtain

$$\mathbb{E} \left[\sum_{t=1}^{T} l_{t}(I_{t}) \right] = \sum_{t=1}^{T} \sum_{i=1}^{N} \mathbb{E} \left[l_{t}(i) \, 1 \{I_{t} = i\} \right] \\
= \sum_{t=1}^{T} \sum_{i=1}^{N} \mathbb{E} \left[l_{t}(i) \right] \mathbb{E} \left[1 \{I_{t} = i\} \right] \\
\left(\text{Since } l_{t}(i) \text{ is independent from } I_{t} \right) \\
= \sum_{t=1}^{T} \sum_{i=1}^{N} m_{i} \mathbb{E} \left[1 \{I_{t} = i\} \right] \\
\left(\text{Since } l_{t}(i) \times V_{i} \text{ and definition of } m_{i} \right) \\
= \sum_{i=1}^{N} m_{i} \mathbb{E} \left[\sum_{t=1}^{T} 1 \{I_{t} = i\} \right] \\
= \sum_{i=1}^{N} m_{i} \mathbb{E} \left[n_{i}(T) \right].$$
Similarly, we easily observe that

$$\min_{1 \leq i \leq N} \mathbb{E} \left[\sum_{t=1}^{T} l_{t}(i) \right] = \mathbb{E} m_{i}^{*}$$

Similarly, we easily observe that

$$\min_{1 \leq i \leq N} \mathbb{E} \left[\sum_{t=1}^{T} l_{t}(i) \right] = Tm^{*}$$

$$= \sum_{i=1}^{N} m^{*} \mathbb{E} \left[n_{i}(T) \right],$$
Since dearly $\sum_{i=1}^{N} n_{i}(T) = T$.

II - Reminder: Sub-gaussian distributions

In the sequel, we'll focus on the case where the loss distributions {vi };=, are subgaussian. This restriction is made for simplicity and can be easily removed (at the price of worse regret bounds).

Definition

A probability distribution V over R is said sub-gaussian with parameter $\sigma > 0$ of Γ , given a random variable $X \times V$, we have

 $\frac{\lambda(X-EX)}{\forall \lambda \in \mathbb{R}} : \log E \in \frac{\lambda^2 \sigma^2}{2}$

Kemanks.

- The terminology sub-gaussian comes from the fact that, if V is the gaussian distribution $\mathcal{N}(m, \sigma^2)$ then, $\forall A \in \mathbb{R}$, $\log E \in \mathbb{R}$ = $\frac{\lambda^2 \sigma^2}{2}$
- When v is a general sub-gaussian distribution with parameter $\sigma^2 > 0$, we'll always implicitly assume that o'is the smallest constant for which the property of the definition holds, i.e. that

$$\sigma^2 = \sup_{\lambda \neq 0} \frac{2}{\lambda^2} \log E e^{\lambda (X - EX)}$$

of is also called the variance proxy of vand it may be shown that inequality

$$Var(v) \leq \sigma^2$$

always holds.

As shown by the next result, the set of sub-gaussian distributions includes all distributions with bounded support.

Lemma 2 (Hoeffoling Lemma)

Let X be a real-valued random variable such that $P(a \le X \le b) = 1$ for some

a < b \in \mathbb{R}. Then, $\forall \lambda \in \mathbb{R}$, \\
\left| \left

In otherwords, the previous lemma states that any probability distribution supported on a finite interval [a, b] is subgaussian with parameter

$$\sigma^2 = \frac{\left(b - a\right)^2}{4}$$

Next, we review the concentration properties of sub-gaussian random variables. From now on, we say that a random variable X is sub-gaussian if its distribution is sub-gaussian.

Suppose X is sub-gaussian with parameter $\sigma^2 > 0$. Then, $\forall t > 0$, $-\frac{t^2}{2\sigma^2}$ $P(X-EX>t) \le e^{-\frac{t^2}{2\sigma^2}}$, $P(X-EX<-t) \le e^{\frac{2}{2\sigma^2}}$ and therefore $\frac{t}{2\sigma^2}$ $P(X-EX|>t) \le e^{\frac{2}{2\sigma^2}}$

Proof: It is enough to prove the first inequality. Indeed, the 2nd follows immediatly by observing that "X sub-gaussian with parameter o²>0" (=> "-X sub-gaussian with parameter o²>0". The 3rd inequality follows also directly by combining the 1st and 2nd inequalities.

Now, observe that $\forall \lambda > 0$, $\lambda(X - E \times) = P(\mathcal{E}) \Rightarrow \mathcal{E}$ $= \lambda t \quad \lambda(X - E \times) \quad (\text{Markov's ineq.})$ $\leq \mathcal{E} \quad \mathcal{E} \quad (\mathcal{E}) \quad$

Lemma 4. Suppose $\{X_i\}_{i=1}^n$ are independent sub-gaussian random variables with parameters $\{T_i\}_{i=1}^2$ respectively. Then

respectively. Then $\sum_{i=1}^{n} X_{i}$ i = 1is sub-gaussian with parameter $\sigma^{2} \leqslant \sum_{i=1}^{2} \sigma_{i}^{2}$.

Proof: Denote $Y_i = X_i - EX_i$. Then, by independence, we get $\lambda \sum_i Y_i$ λY_i

By assumption, we then deduce that

$$\log \mathbb{E} e^{\lambda \sum_{i} Y_{i}^{2}} \leq \sum_{i} \frac{\lambda^{2} \sigma_{i}^{2}}{2} = \frac{\lambda^{2}}{2} \left(\sum_{i} \sigma_{i}^{2} \right),$$

which concludes the proof.

We arrive at the most important fact needed in the sequel, which early follows by combining Lemmas 3 and 4.

Corollary 5

Suppose $X_1,...,X_n$ iid and Mb-gaussians with parameter $\sigma^2 > 0$. Then $\forall \alpha, t > 0$ such that $t^{\alpha} > 1$,

$$P\left(\frac{1}{n}\sum_{i=1}^{n}X_{i}-EX_{1} \geqslant \sigma\sqrt{\frac{2\alpha \ln(t)}{n}}\right)$$
 and
$$P\left(\frac{1}{n}\sum_{i=1}^{n}X_{i}-EX_{1} \leqslant -\sigma\sqrt{\frac{2\alpha \ln(t)}{n}}\right) \leqslant \frac{1}{t^{\alpha}}$$

Proof: According to Lemma 4, $\sum_{i=1}^{n} X_i$ is sub-gaussian with parameter $\leq n \sigma^2$. Hence, Lemma 3 implies that

$$P\left(\frac{1}{n}\sum_{i=1}^{n}X_{i}-EX_{1}\right) \geq P\left(\sum_{i=1}^{n}X_{i}-E\left[\sum_{i=1}^{n}X_{i}\right]\right) \times \exp\left(-\frac{(n\epsilon)^{2}}{2n\sigma^{2}}\right)$$

$$= \exp\left(-\frac{n\epsilon^{2}}{2\sigma^{2}}\right).$$

Setting
$$\frac{1}{t^{\alpha}} = \exp\left(-\frac{n\epsilon^{2}}{2\sigma^{2}}\right)$$
, the previous inequality reads $(\forall x, t > 0 \text{ such that } t^{\alpha} > 1)$

$$P\left(\frac{1}{n}\sum_{i=1}^{n}X_{i} - E[X_{1}] > \sigma\left(\frac{2\alpha\log t}{n}\right) < \frac{1}{t^{\alpha}}\right)$$
Changing X_{i} to $-X_{i}$ implies the second bound.

III Upper Confidence Bound (UCB) Strategy.

In this section, we are back to the stochastic bandit problem introduced in Section I. We'll work under the following assumption:

Assumption (Sub-gaussian losses) The unknown loss distributions $v_1,..., v_N$ are all sub-gaussian with parameter $\sigma^2 > 0$ for a known $\sigma^2 > 0$.

The UCB algorithm we'll describe next combines exploration (of the behavior of the different actions) with exploitation (of good actions already identified)

UCB algorithm

Parameters: Subgaussian parameter $\sigma^2 > 0$ and some parameter $\alpha > 2$.

Initialisation: $n_{i}(0) = 0$, $\hat{m}_{i}(0) = 0$, $\forall i \in [N]$.

For t≯l:

$$\rightarrow$$
 Select $I_t \in argmin B_t(i)$
 $1 \le i \le N$

where
$$\frac{1 \le x \le N}{n_i(t-1)} = \begin{cases}
\hat{m}_i(t-1) - \sigma \sqrt{\frac{2\alpha \ln(t)}{n_i(t-1)}} & \text{if } n_i(t-1) > 0 \\
- \infty & , \text{if } n_i(t-1) = 0
\end{cases}$$
and where
$$\frac{\hat{m}_i(t-1)}{n_i(t-1)} = \frac{1}{n_i(t-1)} \sum_{s=1}^{\infty} l_s(T_s) \, II\{T_s = i\}.$$

$$m_{i}(t-1) = \frac{1}{n_{i}(t-1)} \sum_{s=1}^{t-1} l_{s}(I_{s}) 1\{I_{s}=i\}.$$

is the sample average of all losses obtained from action i at time t-1.

$$\rightarrow$$
 Receive $l_t(I_t)$ and update:
 $n_i(t) := n_i(t-1) + 1\{I_t = i\}$

and
$$n_{i}(t-1) \hat{m}_{i}(t-1) + l_{t}(I_{t}) 1 \{I_{t}=i\}$$
 $n_{i}(t)$

Remark: We have phrased the bandit pbl in terms of "losses", instead of "gains" as is usually the case in the literature, to stay coherent with the previous chapter. In this framework, the term Bt(i) has the flavour of a "lower" bound for a confidence interval of mi instead of an "upper" bound suggested by the name of the algorithm. In the UCB algorithm, the regime switches from a more exploratory phase in the beginning (the confidence intervals are mide in the beginning) to exploitation as time goes on (we come to identify m* with large prob).

Next we study the performance of the UCB algorithm.

Theorem 6

Supposing the distributions of losses are all subgaumian with parameter $\sigma^2 > 0$, the UCB algorithm with parameter $\alpha > 2$ satisfies

$$R_{T} \leqslant \sum_{i:\Delta_{i}>0} \Delta_{i} \left(8\sigma^{2} \frac{\ln T}{\Delta_{i}^{2}} + \frac{\alpha}{\alpha-2}\right)$$

Proof: Recall from Lemma 1 that

$$R_{T} = \sum_{i: \Delta_{i} > 0} \Delta_{i} E \left[n_{i}(T) \right]$$

To prove the result, we therefore need to show that, $\forall i \in [N]$ such that $\Delta_i > 0$, we have the bound

$$\mathbb{E}\left[n_{i}(T)\right] < 8\sigma^{2}\alpha \frac{\ln T}{\Delta_{i}^{2}} + \frac{\alpha}{\alpha - 2}$$

Hence, from now on, we fix $i \in [N]$ such that $\Delta_i > 0$. Denote

$$t^* = \max \{t > 1 : n_i(t) < \lceil 8\sigma^2 x \frac{\ln T}{\Delta_i^2} \rceil \}$$
, where $\lceil \pi \rceil$ denotes the smallest integer larger or equal to π .

Useful observations about t*:

- a) t* is clearly a random variable, due to the randomness of n; (t), t > 1.
- b) By construction, the first N actions of the UCB algorithm takes each of the N possible actions once. In particular, this implies that, \forall $j \in [N]$, \forall t < N, $n_j(t) < 1$. Since [x] > 1 for x > 0, this implies that $t^* > N$. Also, $\forall j \in [N]$, $\forall t > N$, $n_j(t) > 1$.
- c) Finally, note that we clearly have that

$$n_{i}(t^{*}) = \left[8\sigma^{2} \frac{\ln T}{\Delta_{i}^{2}}\right] > 8\sigma^{2} \frac{\ln T}{\Delta_{i}^{2}}$$

In the rest of the proof, we'll look separately at the behaviour of $n_i(T)$ on the event $\{t^*\}T^*\}$ and $\{t^*<T^*\}$.

On the event {t*>T}:

Clearly, we have that

$$E\left[1\left\{t^{*}\right\}T\right]n_{\lambda}(T)\right] \leqslant E\left[1\left\{t^{*}\right\}T\right]\left[8\sigma^{2}\alpha \ln T\right]$$

$$=\left[8\sigma^{2}\alpha \ln T\right]P(t^{*})T$$

$$\leqslant \left(8\sigma^{2}\alpha \ln T + 1\right)P(t^{*})T$$

$$\frac{1}{\Delta_{i}^{2}}$$

On the event 1t* < T :

We can write that

$$E\left[1\left\{t^{*}<\tau\right\} n_{i}(\tau)\right] = E\left[1\left\{t^{*}<\tau\right\} \sum_{t=1}^{N} 1\left\{I_{t}=i\right\}\right]$$

$$= E\left[1\left\{t^{*}<\tau\right\} \left(n_{i}(t^{*}) + \sum_{t=t^{*}+1}^{N} 1\left\{I_{t}=i\right\}\right)\right]$$
Since $n_{i}(t^{*}) = \begin{bmatrix}8\sigma^{2} x \ln T \\ \Delta_{i}^{2}\end{bmatrix} \leqslant \frac{8\sigma^{2} x \ln T}{\Delta_{i}^{2}} + 1$,
we deduce that

$$\mathbb{E}\left[1\left\{t^{*}< T\right\} n_{i}(T)\right]$$

$$\leq \left(\frac{8\sigma^{2}\alpha \ln T}{\Delta_{i}^{2}} + 1\right) \mathbb{P}\left(t^{*}< T\right) + \mathbb{E}\left[1\left\{t^{*}< T\right\}\sum_{t=t^{*}+1}^{1} 1\left\{I_{t}=i\right\}\right]$$

$$+ Combining, * and * we deduce that$$

Combining * and * 2, we deduce that

$$\mathbb{E}\left[n_{i}(T)\right] = \mathbb{E}\left[1\left\{t^{*} > T\right\}n_{i}(T)\right] + \mathbb{E}\left[1\left\{t^{*} < T\right\}n_{i}(T)\right]$$

$$\leq 8 \frac{\sigma^{2} \alpha \ln T}{\Delta_{i}^{2}} + 1 + \mathbb{E}\left[1\left\{t^{*} < T\right\} \sum_{t=t^{*}+1}^{T} 1\left\{T_{t}=i\right\}\right]$$

Since $1 + \frac{2}{\alpha - 2} = \frac{\alpha}{\alpha - 2}$, the proof is complete provided. we show that provided we show that

$$\mathbb{E}\left[1\left\{t^*<\tau\right\}\sum_{t=t^*+1}^{\tau}1\left\{I_{t}=i\right\}\right]\leqslant \frac{2}{\alpha-2}.$$

In the rest of the proof, we therefore fours on this fact. Let us rewrite

$$E\left[1\left\{t^{*}

$$=\sum_{t=1}^{T}P\left(t^{*}< t\leq T, I_{t}=i\right)$$$$

Then, observe that the following fact holds:

Observation We have the inclusion of events

$$\begin{cases} t^* < t < T, \quad I_t = i \end{cases} \subset A_t \cup B_t$$
where
$$A_t := \begin{cases} \hat{m}_{i^*}(t-1) - \sigma & \frac{2\alpha \ln(t)}{n_{i^*}(t-1)} > m_{i^*} \end{cases}$$
and
$$B_t := \begin{cases} \hat{m}_i(t-1) - \sigma & \frac{2\alpha \ln(t)}{n_i(t-1)} < m_i \end{cases}$$

Proof of this fact: By contradiction, suppose that the above inclusion does not hold, i.e., that we may have: $t^* < t \le T$, $T_t = i$ as well as

$$\hat{m}_{i}*(t-1) = \sigma \sqrt{\frac{2\alpha \ln(t)}{n_{i}*(t-1)}} < m_{i}*$$

$$m_i(t-1) - \sigma \sqrt{\frac{2\alpha \ln(t)}{n_i(t-1)}} > m_i.$$

Then, we get (next page)

Given this observation, we can now write that

Finally we'll show that both of these sums are at most $\frac{1}{\alpha-2}$. We prove this fact for the 1st sum only since the 2nd may be treated similarly.

His

For every $t \ge N+1$, $n_j(t-1) \ge 1$. Hence, for $t \ge N+1$,

$$P(A_t) = P\left(\bigcup_{s=1}^{t-1} \{n_i^*(t-1) = s, A_t\}\right)$$

$$\leq \sum_{s=1}^{t-1} \mathbb{P}\left(n_{i}^{*}(t-1)=s, A_{t}\right)$$
 (union bound)

$$P(A_{t}) = P\left(\begin{array}{c} U \\ S = 1 \end{array} \right) \left(\begin{array}{c} t-1 \\ S = 1 \end{array} \right)$$

$$\leq \sum_{s=1}^{t-1} P\left(\begin{array}{c} n_{i}^{*}(t-1) = s, A_{t} \end{array} \right) \quad (union bound)$$

$$\leq \sum_{s=1}^{t-1} P\left(\begin{array}{c} 1 \\ S \\ \overline{j} = 1 \end{array} \right) \left(\begin{array}{c} i^{*} \\ i^{*} \end{array} \right) - \sigma \left(\begin{array}{c} 2\alpha \ln(t) \\ \overline{s} \end{array} \right) \approx m_{i}^{*}$$

$$\leq \sum_{s=1}^{t-1} P\left(\begin{array}{c} 1 \\ S \\ \overline{j} = 1 \end{array} \right) \left(\begin{array}{c} i^{*} \\ S \end{array} \right) - \sigma \left(\begin{array}{c} 2\alpha \ln(t) \\ \overline{s} \end{array} \right) \approx m_{i}^{*}$$

$$\leq \sum_{s=1}^{t-1} P\left(\begin{array}{c} 1 \\ S \\ \overline{j} = 1 \end{array} \right) \left(\begin{array}{c} i^{*} \\ S \end{array} \right) - \sigma \left(\begin{array}{c} 2\alpha \ln(t) \\ \overline{s} \end{array} \right) \approx m_{i}^{*}$$

$$\leq \sum_{s=1}^{t-1} P\left(\begin{array}{c} 1 \\ S \\ \overline{s} = 1 \end{array} \right) \left(\begin{array}{c} \sin \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \sin \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \sin \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \sin \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \sin \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \sin \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \sin \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \sin \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha \\ \cos \alpha \\ \cos \alpha \\ \end{array} \right) \left(\begin{array}{c} \cos \alpha \\ \cos \alpha$$

$$\begin{array}{c|c}
 & t-1 \\
 & \leq \sum_{S=1}^{\infty} \frac{1}{t^{\alpha}} \quad (Corollary 5) \\
 & t = 1 \\
 & t = 1
\end{array}$$

$$= (t-1)t^{-\alpha}$$

This concludes the proof.